

Congestion Control in Wireless Networks

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Motivation

- **Wired Networks**

- Packet loss only due to traffic congestion
- TCP widely used for congestion control

- **Wireless Networks**

- Congestion and Channel Fading both contribute to packet loss
- *Why distinguish?*: Both have different statistics

Thus, TCP needs to be revisited for use in wireless networks

'Control'ing Congestion

- Route Incidence matrix $[[R_{li}]]$, Transmission Rates $x_i(t)$, Link Prices $p_l(t)$

$$y_l(t) = \sum_i R_{li} x_i(t)$$

$$q_i(t) = \sum_l R_{li} p_l(t)$$

- Control Mechanism (State Space formulation)
 - Adapt the rates $x_i(t)$ based on *feedback* $q_i(t)$ (e.g. TCP window control)

$$\dot{z}_i = F_i(z_i, q_i)$$

$$x_i = G_i(z_i, q_i)$$

- The link prices $p_l(t)$ change depending on $y_l(t)$ (e.g. AQM)

$$\dot{v}_l = H_l(v_l, y_l)$$

$$p_l = K_l(v_l, y_l)$$

'Control'ing Congestion (contd...)

- Can be interpreted as a Utility Maximisation Problem *i.e.*

$$\begin{aligned} & \max_{\mathbf{x} \geq \mathbf{0}} \sum_i U_i(x_i) \\ & \text{subject to } R\mathbf{x} \leq \mathbf{C} \end{aligned}$$

where, \mathbf{C} is the link capacity vector

- Different variants of TCP have different utility functions, and different rate and link price control mechanisms
- e.g. TCP Reno

$$U_i(x_i) = \frac{\sqrt{2}}{\tau_i} \tan^{-1}\left(\frac{\tau_i x_i}{\sqrt{2}}\right)$$

Modification for Wireless Networks

- Link Capacities C_l *random* functions of time
- Need to track channel variations and take appropriate action
- A possible approach uses dependence of C on transmit power P
- Thus, modified optimisation problem with additional optimisation variable P

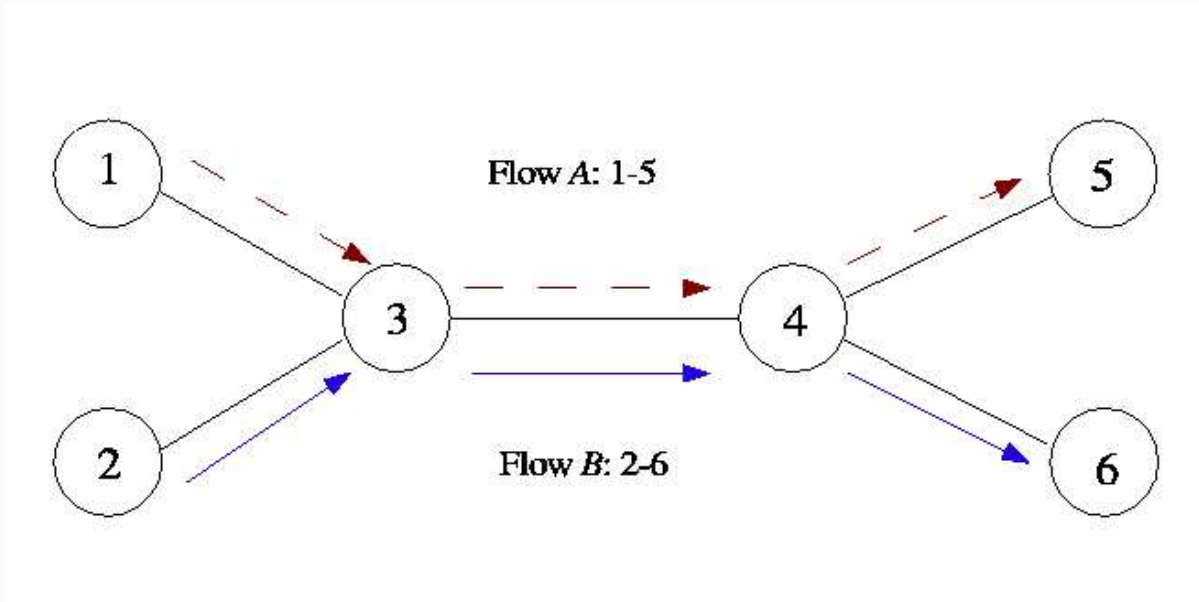
$$\begin{aligned} \max_{\mathbf{x} \geq \mathbf{0}} \quad & \sum_i U_i(x_i) \\ \text{subject to} \quad & R\mathbf{x} \leq \mathbf{C}(\mathbf{P}) \\ & P_l \leq P_{l,max} \quad \forall l \\ & \mathbf{P}, \mathbf{x} \geq \mathbf{0} \end{aligned}$$

- Chiang and Man, break up the above into two separate optimisations

$$\begin{aligned} \text{maximise}_{\mathbf{x} \geq \mathbf{0}} \quad & \sum_s U_s(x_s) - \sum_s \sum_{l \in L(s)} p_l x_s \\ \text{maximise}_{\mathbf{P}_{max} \geq \mathbf{P} \geq \mathbf{0}} \quad & \sum_l p_l C_l(\mathbf{P}) \end{aligned}$$

Simulation Environment

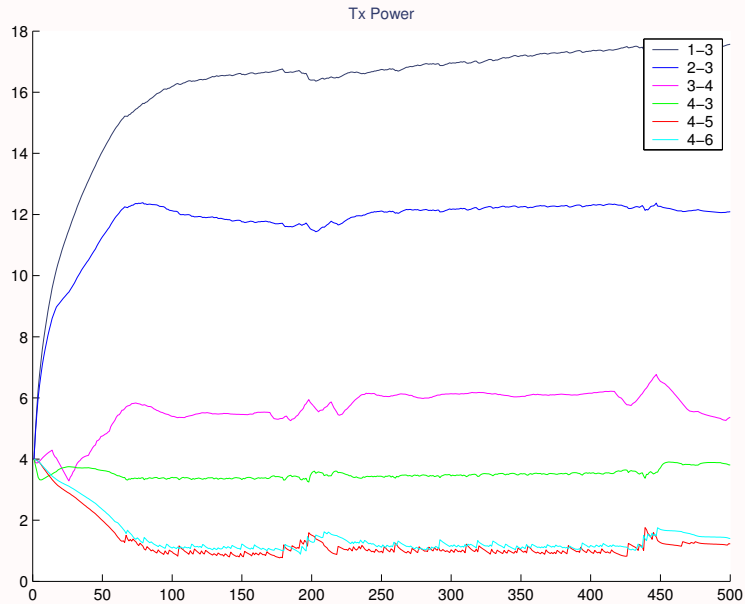
- Topology



- Node: TCP Reno agent whose transmission power can be controlled
- Link: Wireless channel with SIR dependent capacity
- Link: Finite transmission buffer & log-normal channel gain
- Queueing delay and Propagation delays also simulated
- Packet loss either due to buffer overflow (congestion) or due to fading

Results

Start Tx Power	With Power Control	Without Power Control
1	298	220
2	295	228
3	296	221
4	308	233
5	297	162



Results (Contd.)

- Currently experimenting with Network Simulator (NS-2) and trying to understand its functioning and limitations
- Also, looking up more literature on distinguishing between traffic and channel variation statistics
- Will simulate the above formulation and evaluate the performance

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